

Lecture 4: Circuits

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<https://shuaili8.github.io/Teaching/CS445/index.html>

Eulerian circuit

- A closed walk through a graph using every edge once is called an **Eulerian circuit**
- A graph that has such a walk is called an **Eulerian graph**
- **Theorem** (1.2.26, W) A graph G is Eulerian \iff it has at most one nontrivial component and its vertices all have even degree
- (possibly with multiple edges)
- **Proof** “ \implies ” That G must be connected is obvious.
Since the path enters a vertex through some edge and leaves by another edge, it is clear that all degrees must be even

Key lemma

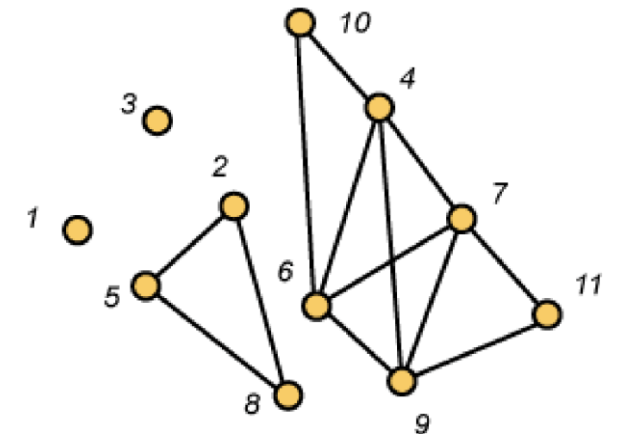
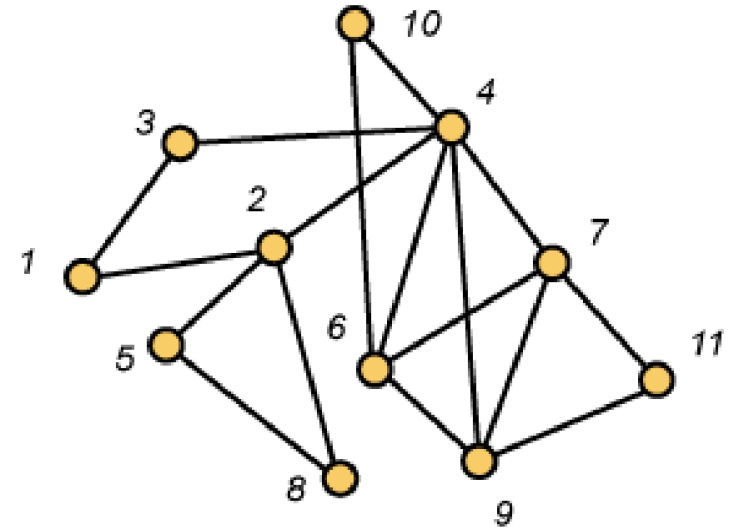
- **Lemma** (1.2.25, W) If every vertex of a graph G has degree at least 2, then G contains a cycle.

Hierholzer's Algorithm for Euler Circuits

1. Choose a root vertex r and start with the trivial partial circuit (r)
2. Given a partial circuit $(x_0, e_1, x_1, \dots, x_{t-1}, e_t, x_t = x_0)$ that traverses not all edges of G , remove these edges from G
3. Let i be the least integer for which x_i is incident with one of the remaining edges
4. Form a greedy partial circuit among the remaining edges of the form $(x_i = y_0, e'_1, y_1, \dots, y_{s-1}, e'_s, y_s = x_i)$
5. Expand the original circuit by setting $(x_0, e_1, \dots, e_i, x_i = y_0, e'_1, y_1, \dots, y_{s-1}, e'_s, y_s = x_i, e_{i+1}, \dots, e_t, x_t = x_0)$
6. Repeat step 2-5

Example

1. Start with the trivial circuit (1)
2. Greedy algorithm yields the partial circuit (1,2,4,3,1)
3. Remove these edges
4. The first vertex incident with remaining edges is 2
5. Greedy algorithms yields (2,5,8,2)
6. Expanding (1,2,5,8,2,4,3,1)
7. Remove these edges



Example (cont.)

6. Expanding (1,2,5,8,2,4,3,1)

7. Remove these edges

8. First vertex incident with remaining edges is 4

9. Greedy algorithm yields (4,6,7,4,9,6,10,4)

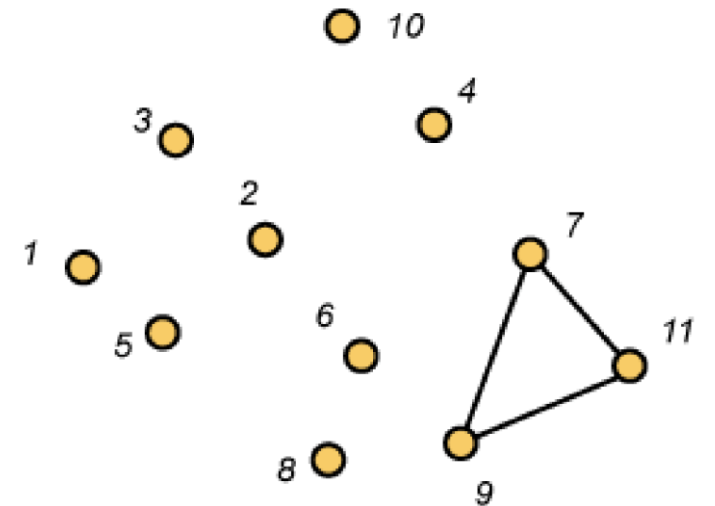
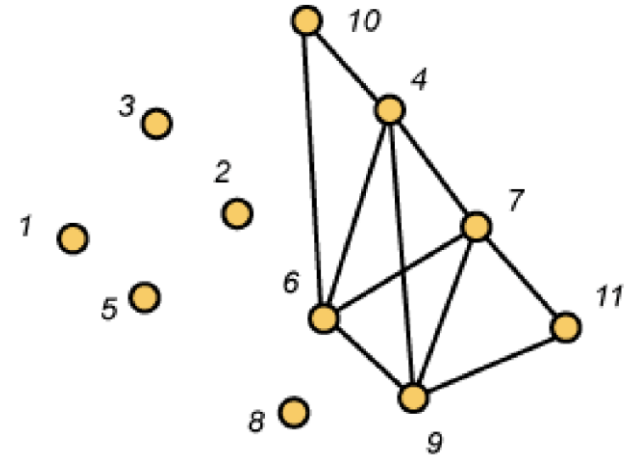
10. Expanding (1,2,5,8,2,4,6,7,4,9,6,10,4,3,1)

11. Remove these edges

12. First vertex incident with remaining edges is 7

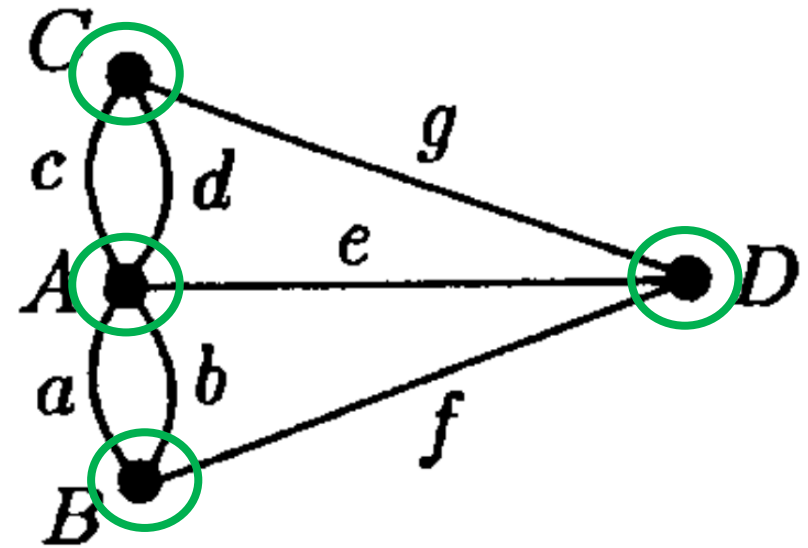
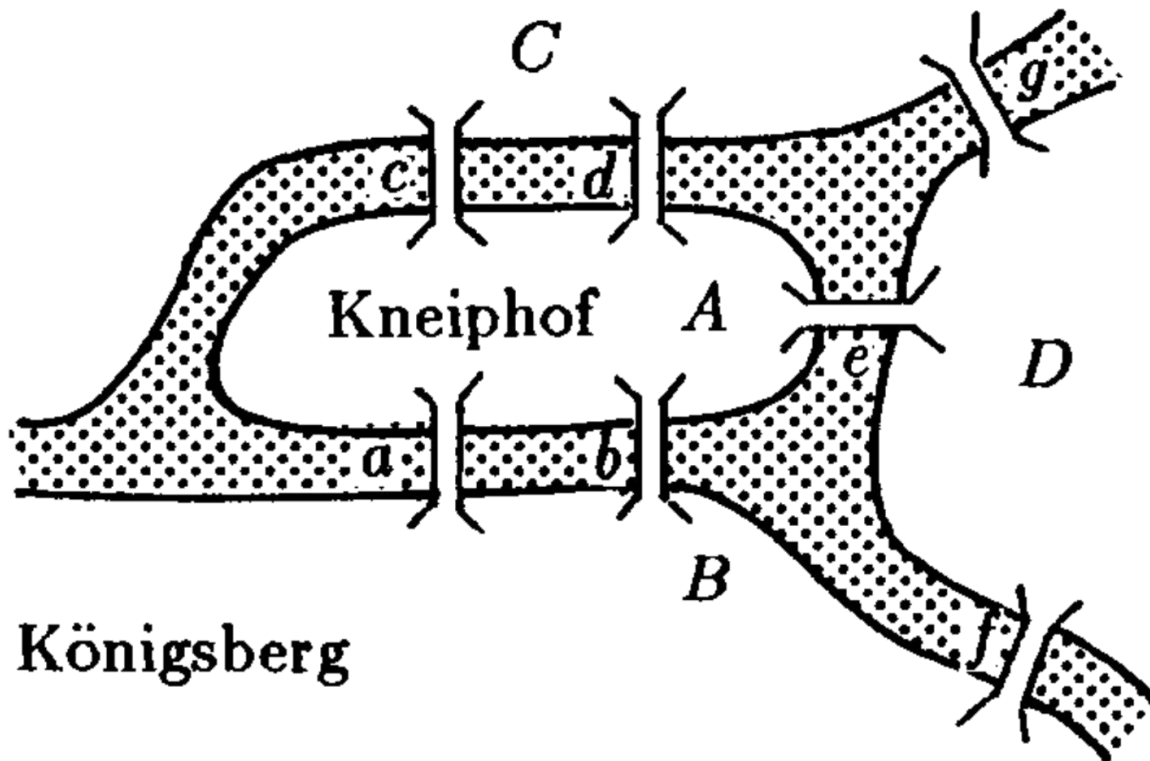
13. Greedy algorithm yields (7,9,11,7)

14. Expanding (1,2,5,8,2,4,6,7,9,11,7,4,9,6,10,4,3,1)



Eulerian circuit

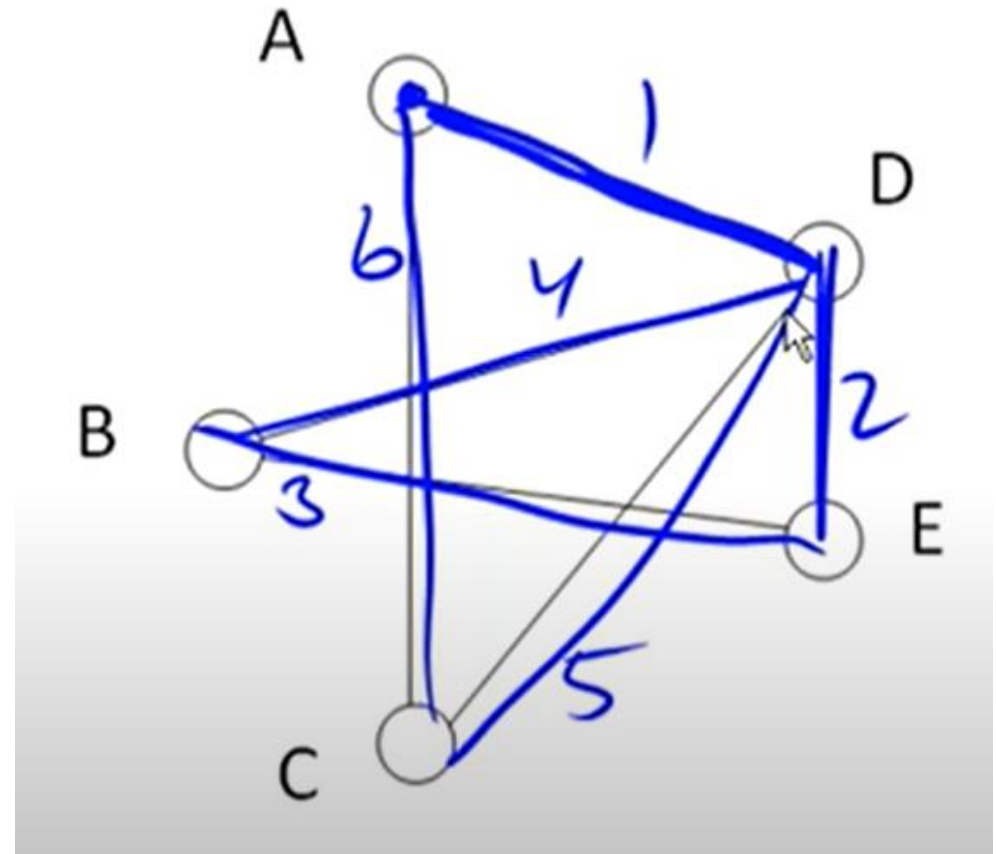
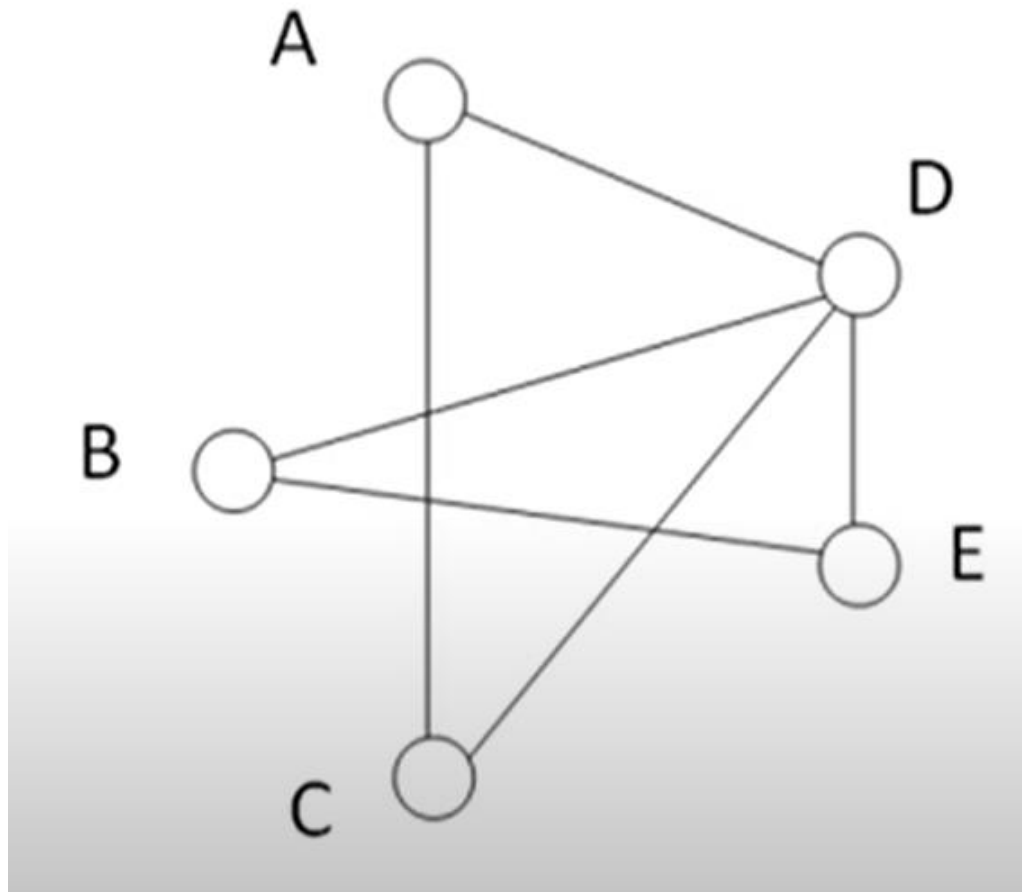
- **Theorem** (1.2.26, W) A graph G is Eulerian \Leftrightarrow it has at most one nontrivial component and its vertices all have even degree



Fleury's Algorithm for Identifying Eulerian Circuits

- (Ex3, S1.4.2, H)
- Given: An Eulerian graph G , with all of its edges *unmarked*
 1. Choose a vertex v , and call it the “lead vertex”
 2. If all edges of G have been marked, then stop. Otherwise continue to step 3
 3. Among all edges incident with the lead vertex, choose, if possible, one that is not a bridge of the subgraph formed by the unmarked edges. If this is not possible, choose any edge incident with the lead. Mark this edge and let its other end vertex be the new lead vertex
 4. Go to step 2

Example



Other properties

- **Proposition** (1.2.27, W) Every even graph decomposes into cycles
- The necessary and sufficient condition for a **directed Eulerian circuit** is that the graph is connected and that each vertex has the same 'in-degree' as 'out-degree'

TONCAS

- **TONCAS**: The obvious necessary condition is also sufficient
- **Theorem** (1.2.26, W) A graph G is Eulerian \Leftrightarrow it has at most one nontrivial component and its vertices all have even degree
- **Proposition** (1.3.28, W) The nonnegative integers d_1, \dots, d_n are the vertex degrees of some graph $\Leftrightarrow \sum_{i=1}^n d_i$ is even
- (Possibly with loops)
- Otherwise $(2,0,0)$ is not realizable

1.3.63. (!) Let d_1, \dots, d_n be integers such that $d_1 \geq \dots \geq d_n \geq 0$. Prove that there is a loopless graph (multiple edges allowed) with degree sequence d_1, \dots, d_n if and only if $\sum d_i$ is even and $d_1 \leq d_2 + \dots + d_n$. (Hakimi [1962])

Hamiltonian path/circuits

- A **path** P is **Hamiltonian** if $V(P) = V(G)$
 - Any graph contains a Hamiltonian path is called **traceable**
- A **cycle** C is called **Hamiltonian** if it spans all vertices of G
 - A graph is called **Hamiltonian** if it contains a Hamiltonian circuit
- In the mid-19th century, Sir William Rowan Hamilton tried to popularize the exercise of finding such a closed path in the graph of the dodecahedron

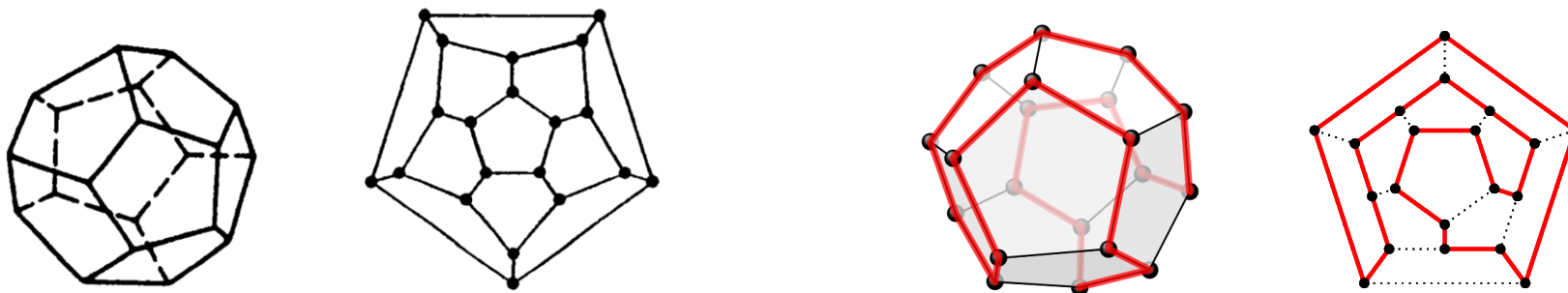
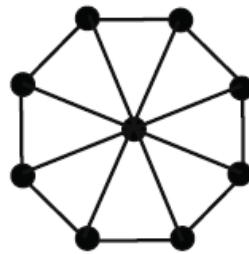


Figure 1.9

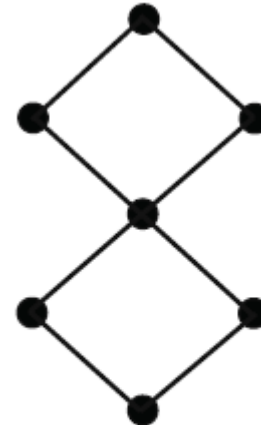
Degree parity is not a criterion

Theorem (1.2.26, W) A graph G is Eulerian \Leftrightarrow it has at most one nontrivial component and its vertices all have even degree

- Hamiltonian graphs
 - all even degrees C_{10}
 - all odd degrees K_{10}
 - a mixture G_1
- non-Hamiltonian graphs
 - all even G_2
 - all odd $K_{5,7}$
 - mixed P_9



G_1



G_2

Example

- The Petersen graph has a Hamiltonian path but no Hamiltonian cycle

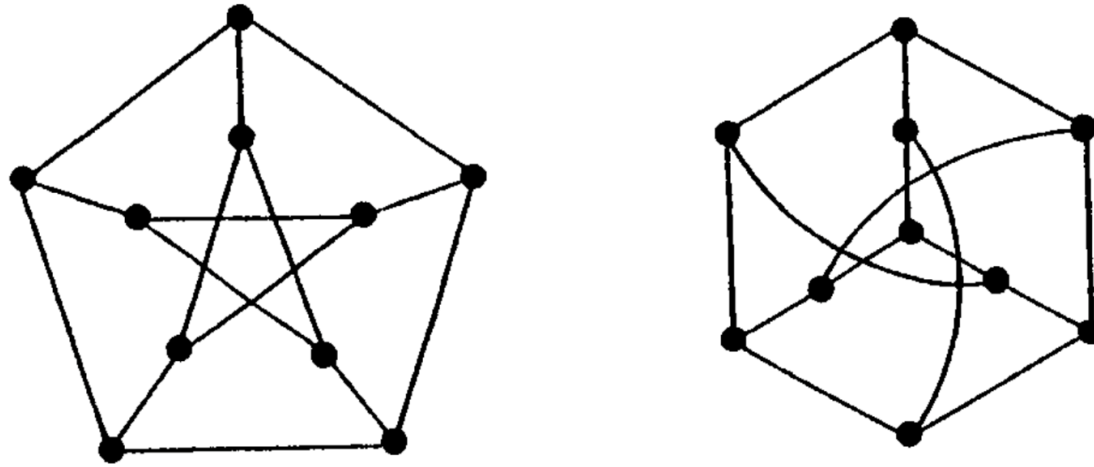
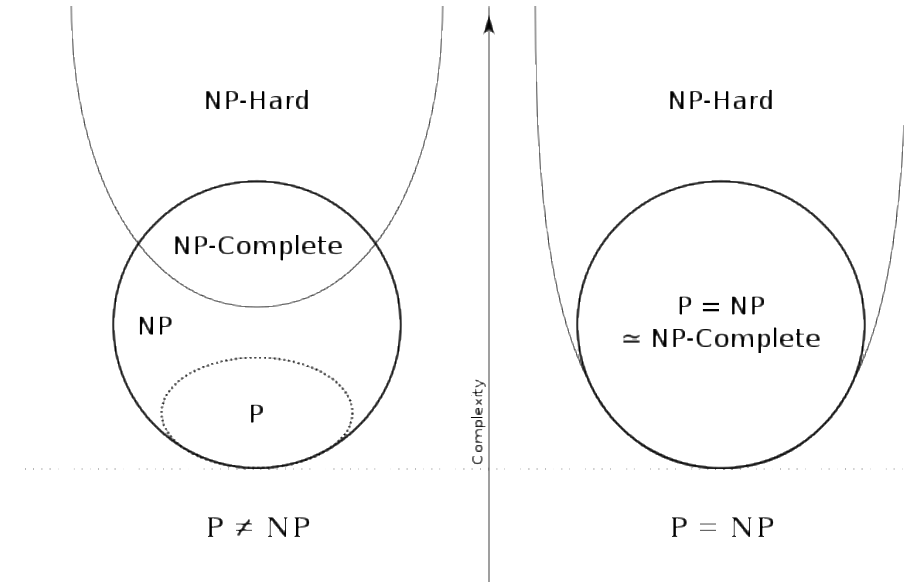


Figure 1.4

- Determining whether such paths and cycles exist in graphs is the Hamiltonian path problem, which is **NP-complete**

P, NP, NPC, NP-hard

- **P** The general class of questions for which some algorithm can provide an answer in polynomial time
- **NP** The class of questions for which an answer can be *verified* in polynomial time
- **NP-Complete**
 1. c is in NP
 2. Every problem in NP is reducible to c in polynomial time
- **NP-hard**
 - ~~c is in NP~~
 - Every problem in NP is reducible to c in polynomial time



Large minimal degree implies Hamiltonian

- **Theorem** (1.22, H, Dirac) Let G be a graph of order $n \geq 3$. If $\delta(G) \geq n/2$, then G is Hamiltonian

Proposition (1.3.15, W) If $\delta(G) \geq \frac{n-1}{2}$, then G is connected

(Ex16, S1.1.2, H) (1.3.16, W)

If $\delta(G) \geq \frac{n-2}{2}$, then G need not be connected

- The bound is **tight**
(Ex12b, S1.4.3, H) $G = K_{r,r+1}$ is not Hamiltonian
- The condition is not necessary
 - C_n is Hamiltonian but with small minimum (and even maximum) degree

Generalized version

- **Exercise** (Theorem 1.23, H, Ore; Ex3, S1.4.3, H) Let G be a graph of order $n \geq 3$. If $\deg(x) + \deg(y) \geq n$ for all pairs of nonadjacent vertices x, y , then G is Hamiltonian

Theorem (1.22, H, Dirac) Let G be a graph of order $n \geq 3$. If $\delta(G) \geq n/2$, then G is Hamiltonian

Independence number & Hamiltonian

- A set of vertices in a graph is called **independent** if they are pairwise nonadjacent
- The **independence number** of a graph G , denoted as $\alpha(G)$, is the largest size of an independent set
- Example: $\alpha(G_1) = 2, \alpha(G_2) = 3$
- **Theorem** (1.24, H) Let G be a connected graph of order $n \geq 3$. If $\kappa(G) \geq \alpha(G)$, then G is Hamiltonian
- The result is **tight**: $\kappa(G) \geq \alpha(G) - 1$ is not enough
 - $K_{r,r+1}: \kappa = r, \alpha = r + 1$
 - Peterson graph: $\kappa = 3, \alpha = 4$ (Ex4, S1.4.3, H)

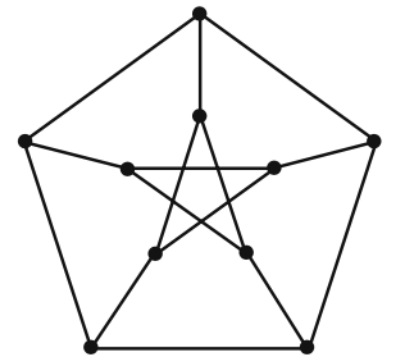
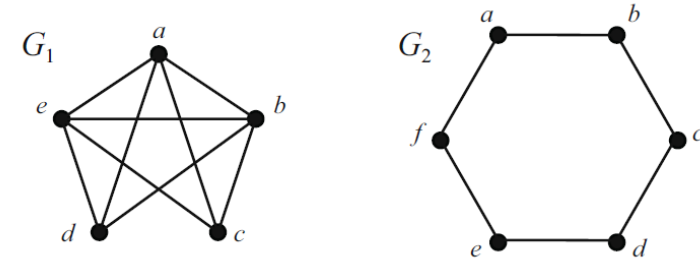
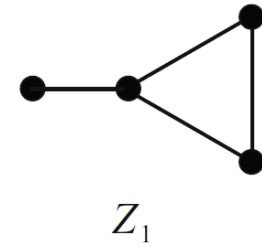
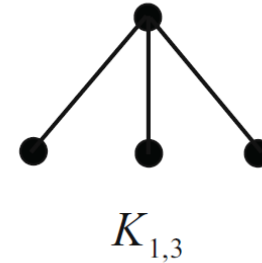


FIGURE 1.63. The Petersen Graph.

Pattern-free & Hamiltonian



- G is H -free if G doesn't contain a copy of H as induced subgraph
- **Theorem** (1.25, H) If G is 2-connected and $\{K_{1,3}, Z_1\}$ -free, then G is Hamiltonian

(Ex14, S1.1.2, H) $\kappa(G) \geq 2$ implies G has at least one cycle

- The condition 2-connectivity is necessary
- (Ex2, S1.4.3, H) If G is Hamiltonian, then G is 2-connected