

# Lecture 7: Planarity

Shuai Li

John Hopcroft Center, Shanghai Jiao Tong University

<https://shuaili8.github.io>

<https://shuaili8.github.io/Teaching/CS445/index.html>

# Motivation

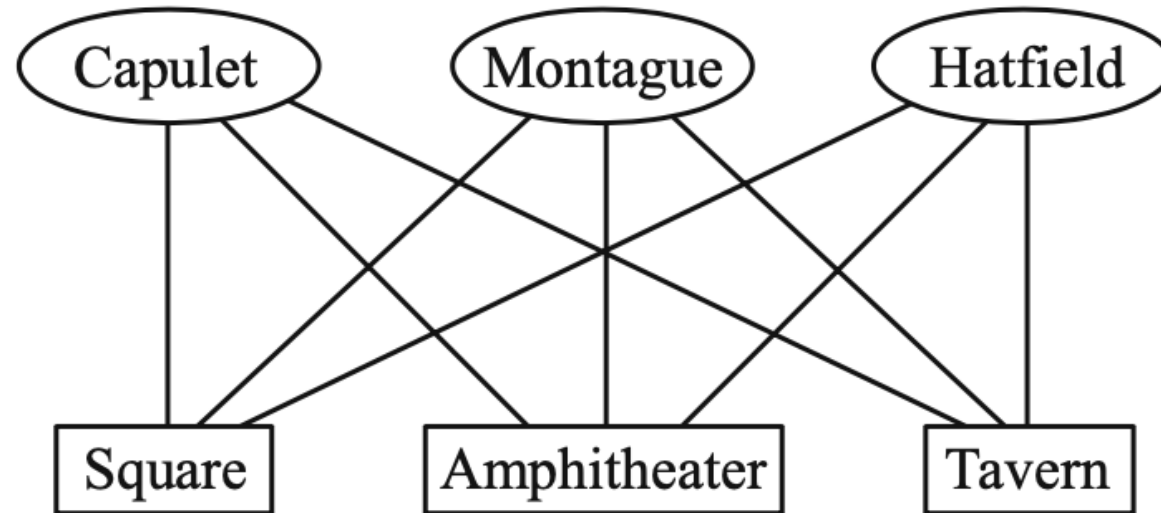


FIGURE 1.72. Original routes.

# Definition and examples

- A graph  $G$  is said to be **planar** if it can be drawn in the plane in such a way that pairs of edges intersect only at vertices
- If  $G$  has no such representation,  $G$  is called **nonplanar**
- A drawing of a planar graph  $G$  in the plane in which edges intersect only at vertices is called a **planar representation** (or a planar embedding) of  $G$

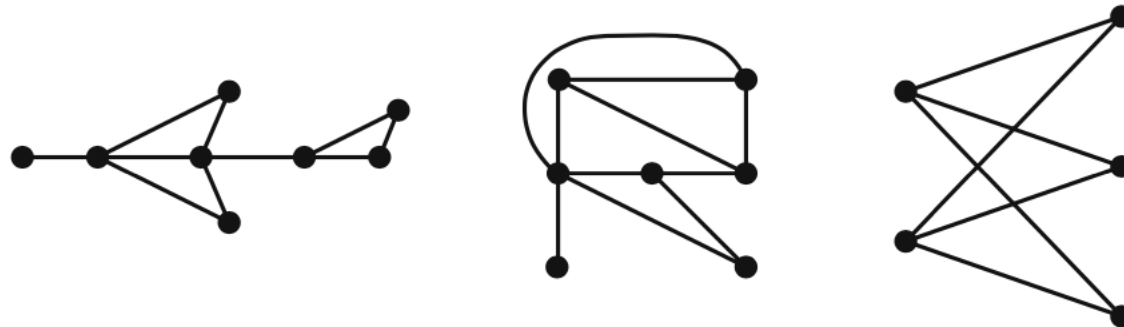
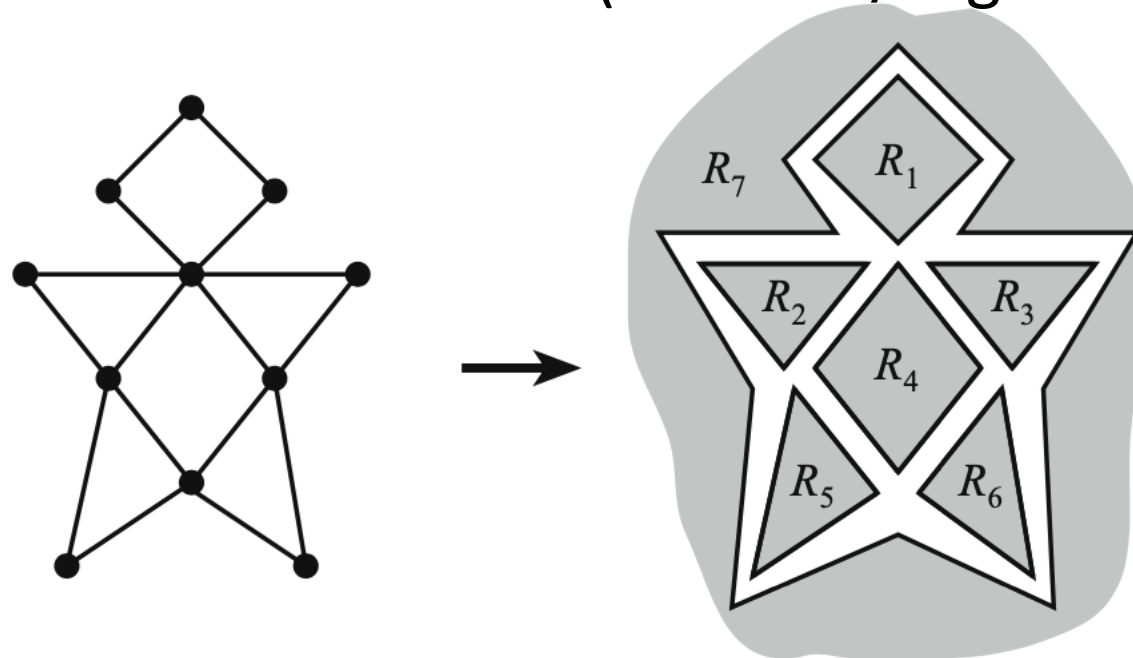


FIGURE 1.73. Examples of planar graphs.

# Region

- Given a planar representation of a graph  $G$ , a **region** is a maximal section of the plane in which any two points can be joined by a curve that does not intersect any part of  $G$
- The region  $R_7$  is called the **exterior** (or outer) region



# An edge bounds a region

- An edge can come into **contact** with either one or two regions
- Example:
  - Edge  $e_1$  is only in contact with one region  $S_1$
  - Edge  $e_2, e_3$  are only in contact with  $S_2$
  - Each of other edges is in contact with two regions
- An edge  $e$  **bounds** a region  $R$  if  $e$  comes into contact with  $R$  and with a region **different** from  $R$
- The **bounded degree**  $b(R)$  is the number of edges that bound the region
  - Example:  $b(S_1) = b(S_3) = 3, b(S_2) = 6$

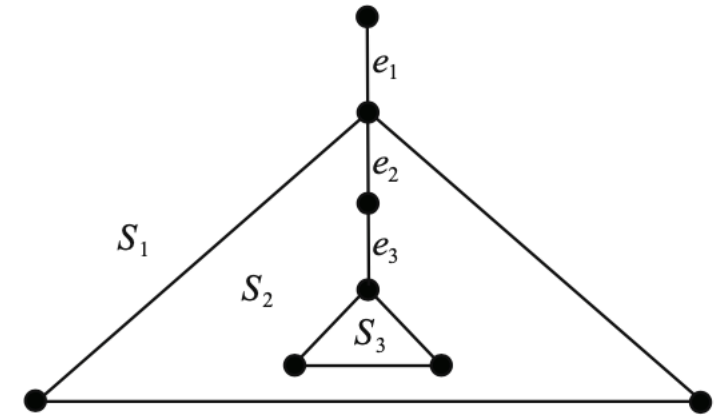
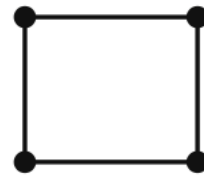


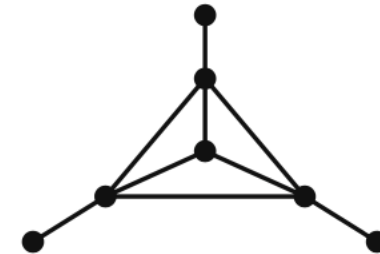
FIGURE 1.76. Edges  $e_1, e_2,$  and  $e_3$  touch one region only.

# The relationship between numbers of vertices, edges and regions

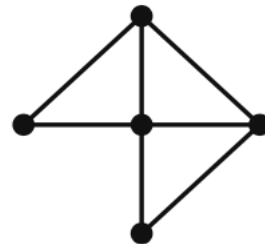
- The number of vertices  $n$
- The number of edges  $m$
- The number of regions  $r$



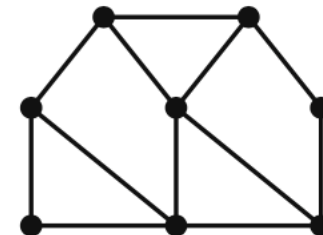
$n = 4$   
 $m = 4$   
 $r = 2$



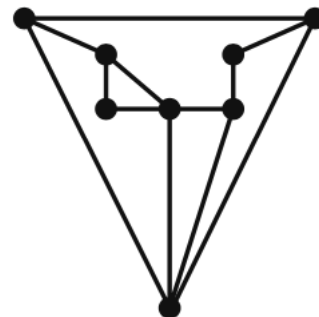
$n = 7$   
 $m = 9$   
 $r = 4$



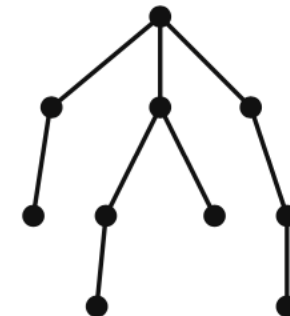
$n = 5$   
 $m = 7$   
 $r = 4$



$n = 8$   
 $m = 12$   
 $r = 6$



$n = 8$   
 $m = 12$   
 $r = 6$



$n = 10$   
 $m = 9$   
 $r = 1$

# Euler's formula

- **Theorem** (1.31, H; Euler 1748) If  $G$  is a connected planar graph with  $n$  vertices,  $m$  edges, and  $r$  regions, then

$$n - m + r = 2$$

- Need Lemma: (Ex4, S1.5.1, H) Every tree is planar
- (Ex6, S1.5.2, H) Let  $G$  be a planar graph with  $k$  components. Then

$$n - m + r = k + 1$$

$K_{3,3}$  is nonplanar

- **Theorem** (1.32, H)  $K_{3,3}$  is nonplanar

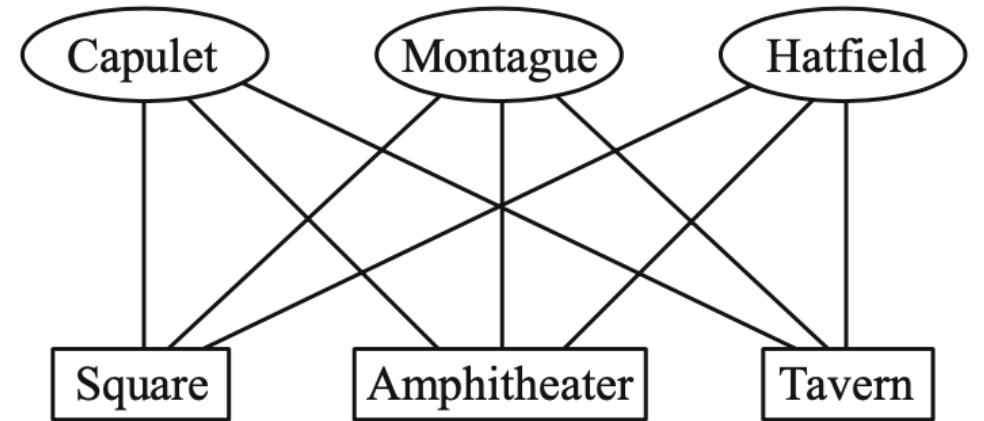


FIGURE 1.72. Original routes.



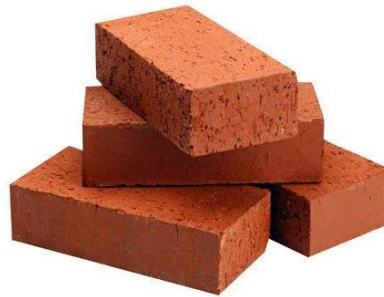
# Upper bound for $m$

- **Theorem** (1.33, H) If  $G$  is a planar graph with  $n \geq 3$  vertices and  $m$  edges, then  $m \leq 3n - 6$ . Furthermore, if equality holds, then every region is bounded by 3 edges.
- (Ex4, S1.5.2, H) Let  $G$  be a connected, planar,  $K_3$ -free graph of order  $n \geq 3$ . Then  $G$  has no more than  $2n - 4$  edges
- **Corollary** (1.34, H)  $K_5$  is nonplanar
- **Theorem** (1.35, H) If  $G$  is a planar graph, then  $\delta(G) \leq 5$
- (Ex5, S1.5.2, H) If  $G$  is bipartite planar graph, then  $\delta(G) < 4$

# Polyhedra

# (Convex) Polyhedra 多面体

- A **polyhedron** is a solid that is bounded by flat surfaces



# Polyhedra are planar

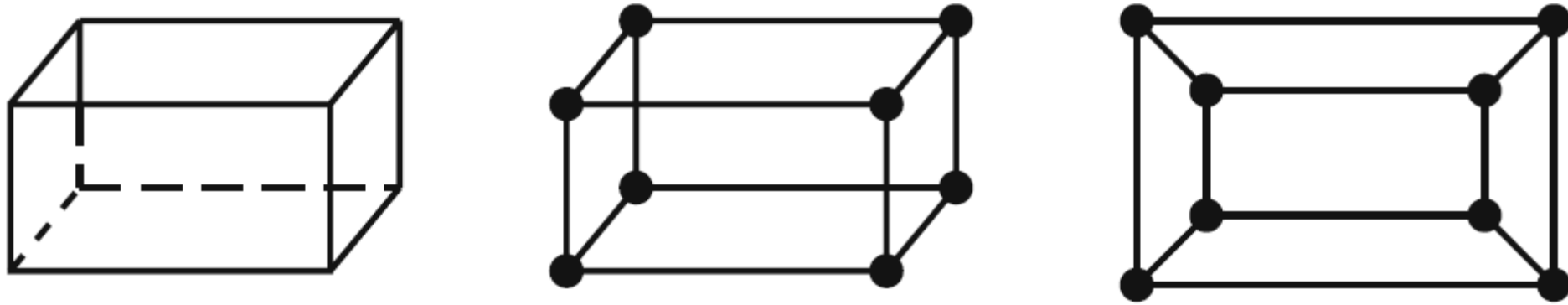


FIGURE 1.81. A polyhedron and its graph.

# Properties

- **Theorem** (1.36, H) If a polyhedron has  $V$  vertices,  $E$  edges, and  $F$  faces, then

$$V - E + F = 2$$

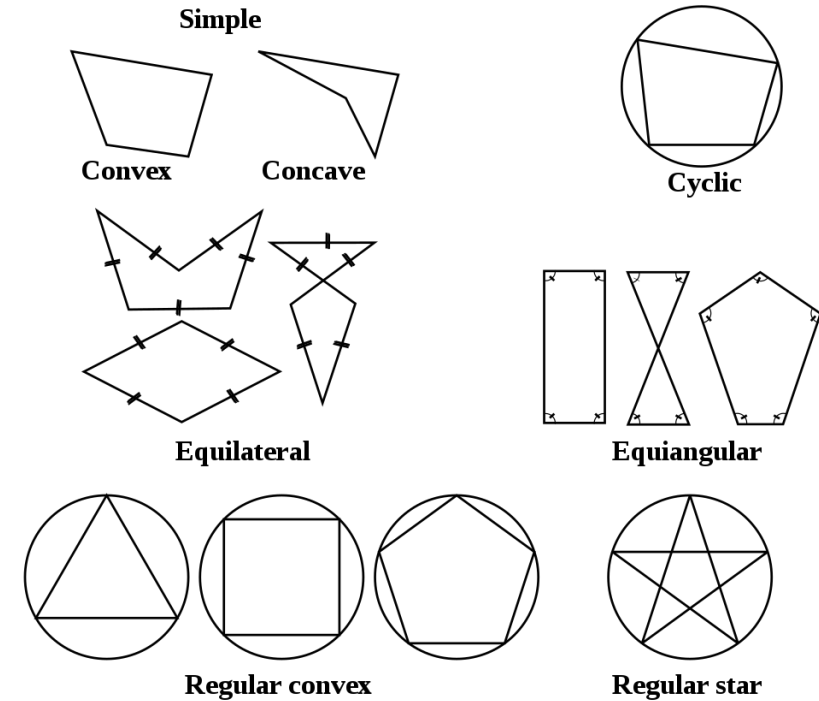
- Given a polyhedron  $P$ , define

$$\rho(P) = \min\{b(R) : R \text{ is a region of } P\}$$

- **Theorem** (1.37, H) For all polyhedron  $P$ ,  $3 \leq \rho(P) \leq 5$

# Regular polyhedron 正多面体

- A **regular polygon** is one that is equilateral and equiangular  
正多边形(cycle), 等边、等角
- A polyhedron is **regular** if its faces are mutually congruent, regular polygons and if the number of faces meeting at a vertex is the same for every vertex  
正多面体  
面是相互全等的、正多边形、点的度数相等



# Regular polyhedron 正多面体

- **Theorem** (1.38, H) There are exactly five regular polyhedral
- 正四面体
- 立方体（正六面体）
- 正八面体
- 正十二面体
- 正二十面体

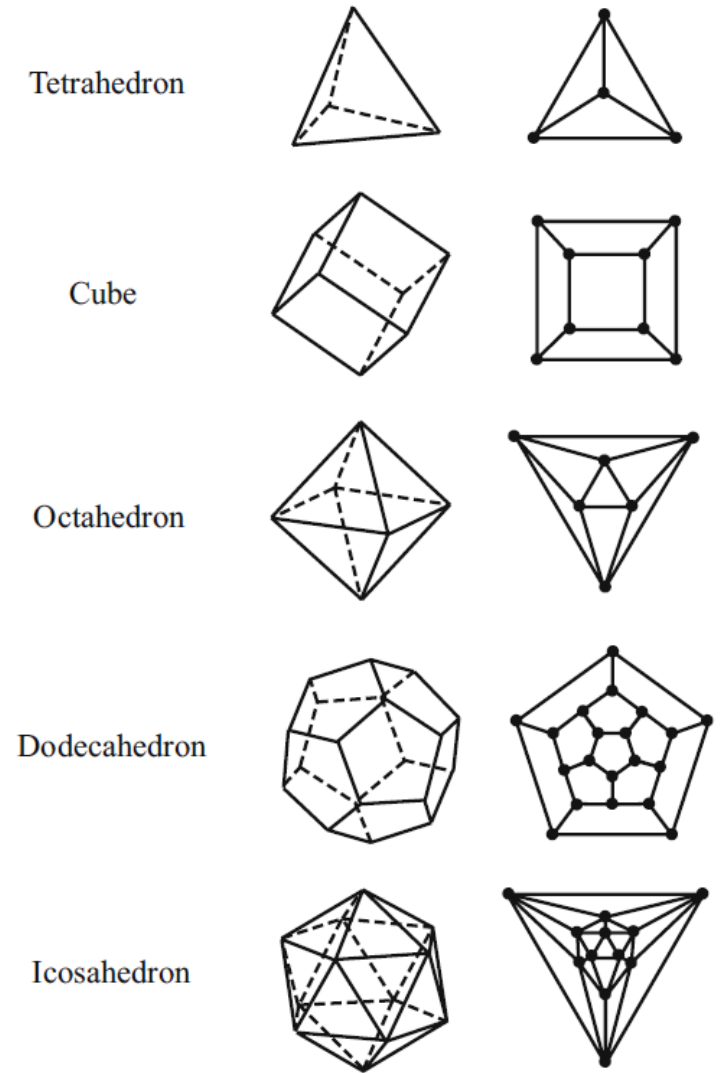


FIGURE 1.82. The five regular polyhedra and their graphical representations.

# Kuratowski's Theorem



# Subdivision 细分

- A **subdivision** of an edge  $e$  in  $G$  is a substitution of a path for  $e$
- A graph  $H$  is a **subdivision** of  $G$  if  $H$  can be obtained from  $G$  by a finite sequence of subdivisions
- Example:
  - The graph on the right contains a subdivision of  $K_5$
  - In the below,  $H$  is a subdivision of  $G$

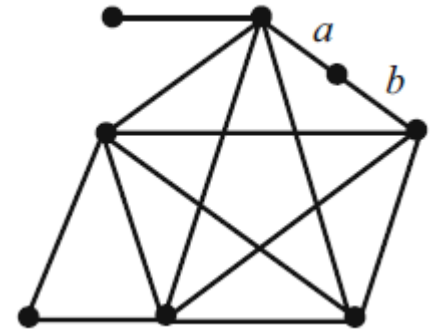
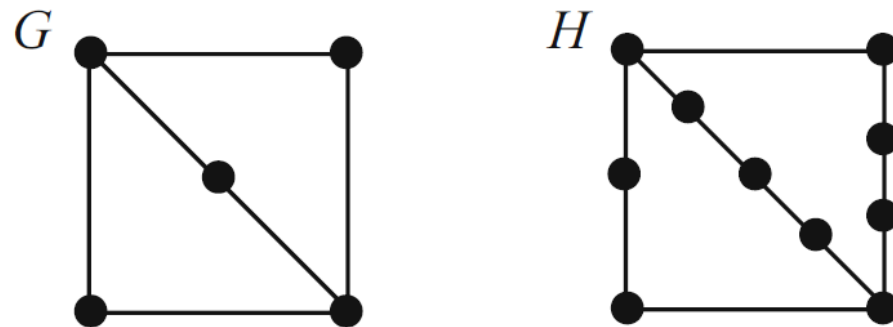


FIGURE 1.84. A graph and a subdivision.

# Kuratowski's Theorem

- **Theorem** (1.39, H; Ex1, S1.5.4, H) A graph  $G$  is planar  $\Leftrightarrow$  every subdivision of  $G$  is planar
- **Theorem** (1.40, H; Kuratowski 1930) A graph is planar  $\Leftrightarrow$  it contains no subdivision of  $K_{3,3}$  or  $K_5$